Formal software development methods

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Lecture #1: Introduction
Motivation

Software validity is one of the main open problems in computer science.

- Bugs have been there forever.  
  You can’t wish them away  
  (like pointer-arithmetic or goto-s)
- Software verification is almost entirely  
  fuelled by ideas from academia.  
  - Theory makes so much sense  
  - Heuristics need theoretical basis to work.  
    Blind heuristics simply don’t work.
Some bugs

- Ariane V Crash (’96)
  64bit -> 16 bit conversion
- Pentium FDIV bug (’97)
  lookup table had mistakes
- Mars Orbiter
  feet-per-second -> Newtons-per-second
- Therac-25
  Radiation therapy machine overdoses patients
Blue screen of death!

Presentation of a Windows 98 beta by Bill Gates at COMDEX on April 20, 1998.
Some bugs...

- In 2006, Windows crashed while Gates was giving a demo of Microsoft Media Center at CES ---

  “blue screen of death”

- Windows 2000 is the chosen OS for the new Type 45 Destroyers of the British Navy.

- Plan to fit it to Vanguard class boats that carry the UK's Trident thermo-nuclear intercontinental ballistic missiles.

- “Windows for Warships safe for Royal Navy” --- says MoD
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Goal

- Make software reliable
- Software industry---Automobile industry
- What is the notion of certification for software? Compare to airline industry.
  In the future, will software be required to meet certain requirements (like being certified by a model-checker) before being used for certain applications?
Goal

- What will JAVA 2010 look like?
  - More features for program annotation
  - Built-in tools that can validate/model-check properties and give feedback to programmer
  - Programming styles that can help automate verification
Review

• ’70s -- proving programs correct
  • Hoare, Djikstra
  • Program invariants
  • Belief: programmers will eventually write
    programs and prove them correct with
    help of theorem provers.

• Failure
  • No way to find bugs; heavily manual; unaccepted in industry
    (and this class!)
  • Fails for large systems
Review

SPIN (Holzmann, Bell Labs, ’90s)

- Explicit-state model checker
- Heuristics to control state-space explosion
  - Partial order reduction
  - Hashing and approximate search
- Specification: LTL / automata
Review

Focus on hardware verification
SMV (McMillan, Clarke, CMU, ’80s)

- Symbolic model checker using binary decision diagrams (BDDs)

- Could handle large state spaces
  - Heuristics to handle search spaces well
  - Specification: CTL (and later LTL)
  - by far the most useful technique in the hardware domain
Review

- Advent of SAT tools (2000)
  - zChaff (Princeton)
  - can handle formulas with 100000 vars, and millions of clauses!

- The idea of bounded model checking
  - Is there a path of length k that reaches an unsafe state?
  - NuSMV and contemporary model checkers use SAT heavily.
Review

- The SLAM tool from Microsoft Research
  Ball and Rajamani, 2000
- Static Driver Verifier (SDV) – distributed tool from Microsoft Research – big breakthrough!

- Model-checker that validates device drivers against formal specifications.

- Key ideas
  - Predicate abstraction to boolean programs
  - Algorithms to check pushdown automata
  - State-space exploration of bool pgms using BDDs.
  - Tremendous success; Static Driver Verifier
Review

- The Metal project from Stanford.
  Dawson Engler

- Static analysis to find patterns of bad programming practice in systems code.

- Very successful in terms of errors found
  - 100s of bugs (incl security) found in Linux/BSD
  - Errors in various protocols, drivers.
  - Coverity (2004)
Review

- Light-weight static analysis
  - Pointer chasing
  - Data-dependency analysis
  - Usually explicit-state analysis on CFG

- Example: Codesurfer (Tom Reps)
Techniques you will learn

- Abstraction into models
  - Abstract Interpretation (Cousot&Cousot ’70s)
  - Predicate abstraction
  - Automatic theorem provers
  - Inference of invariants
  - Shape analysis for handling data structures
  - Separation logic
  - Abstraction refinement
    - Counter-example guided
    - SAT proof guided
Techniques you will learn

- Specification of properties
  - Temporal logics (LTL, CTL), assert, pre-post conditions (FOL)
  - Automata-based techniques to verify LTL
    (Gödel-award winning work!)

- Algorithms to search models
  - Static analysis: Dataflow problems using finite models
  - Finite state models (large!)
    Boolean Decision Diagram (BDDs)
  - Bounded model checking
    SAT techniques
  - Finite-data programs with recursion
    Algorithms to explore models with recursion
Techniques (you will learn)

- Software verification (infinite domains)
  - Floyd’s theory of program verification
    -- reducing whole program verification to normal mathematical theorems
  - Hoare’s pre-post condition verification
  - Automatic theorem provers pave the way for some automated verification of software (use of tool Boogie or variant).

- Automatic software verification using predicate abstraction
Tools you will learn to use

- SMV, NuSMV (originally from CMU)
  - Symbolic Model Verifier; LTL; CTL; BDDs; SAT solving

- Boogie (or JACK)
  - Invariant based verification using theorem proving

- SAT solver (zCHAFF or minsat)
Gritty details

- Undergrad requirement will be simpler and more straightforward
- Graduate/PhD students will have to do a project near the end of the course
- No midterm exam; there will be a final exam
- Homework will be given every 2 weeks
  - Homework will consist of pen-paper questions as well as implementation of program models and use of verification tools
- Grading
  - Undergrads: HW: 40% Finals: 60%
  - Grads/PhD: HW 30%; Project 30%; Finals: 60%
- More info on webpage coming soon.....