

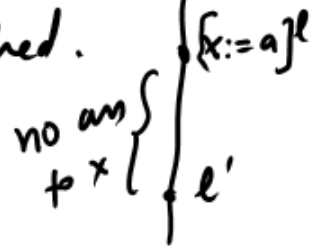
Dataflow Analysis

- Feb 19

```
y := x; ①  
z := 1 ②  
while [y > 1] ③ do (  
    z = z * y ④  
    y := y - 1 ⑤  
)  
y := 0 ⑥
```

Reaching Definitions

An assignment $[x := a]^l$ reaches entry or exit of l' if there is an execution where $[x := a]^l$ occurs and after that there is no assignment to x till l' is reached.



RD is undecidable.

$[x := a]^l ; S ; [y := x]^{l'}$

Programming Language

$a ::= x \mid n \mid a_1 \text{ op}_a a_2$

$b ::= \text{true} \mid \text{false} \mid a_1 \text{ op}_r a_2 \mid$
 $b_1 \text{ op}_b b_2 \mid \text{op}'_b b$

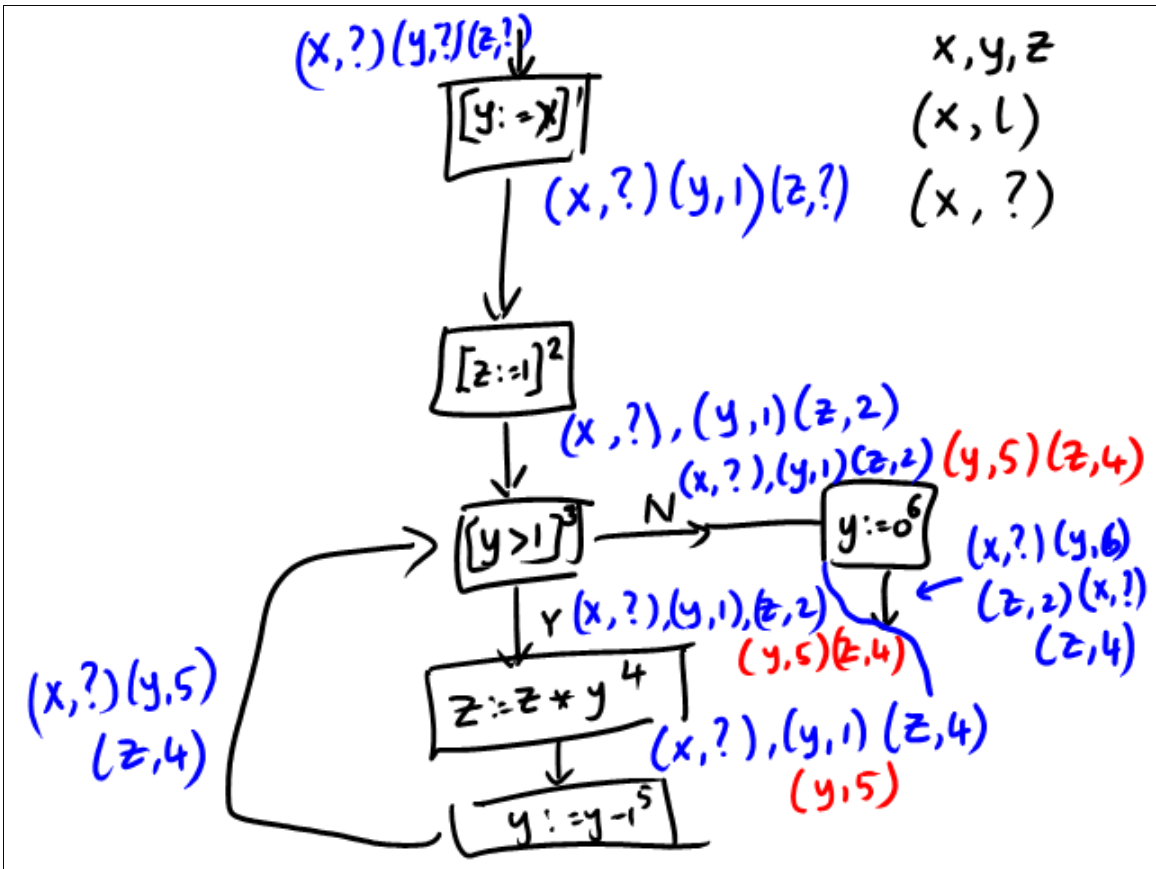
Eg. $\text{op}_a = \{ +, -, \dots \}$

$\text{op}_r = \{ <, =, >, \leq, \dots \}$

$\text{op}_b = \{ \wedge, \vee, \neg \}$

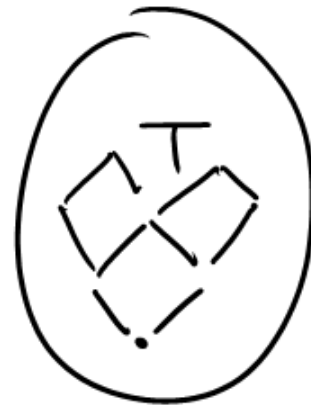
Statements

$S ::= [x := a]^l \mid [skip]^l \mid$
 $\text{if } [b]^l \text{ then } S_1 \text{ else } S_2 \mid$
 $\text{while } [b]^l \text{ do } S \mid$
 $S_1 ; S_2$



In general,

$$\text{if } f : T \rightarrow T$$



RD analysis

$$f : T \rightarrow T$$

$$T = \langle DF_{\text{entry}}', DF_{\text{exit}}', \dots \rangle$$

$$T = \langle DF_{\text{entry}}^i, DF_{\text{exit}}^i \rangle_{i=1, \dots, 6}$$

Lattice:

$$\langle DF_{\text{entry}}^i, DF_{\text{exit}}^i \rangle \leq \langle DF_{\text{entry}}^{i'}, DF_{\text{exit}}^{i'} \rangle$$

$$\text{iff } \forall i. \begin{array}{l} DF_{\text{entry}}^i \subseteq DF_{\text{entry}}^{i'} \\ DF_{\text{exit}}^i \subseteq DF_{\text{exit}}^{i'} \end{array}$$

$f : T \rightarrow T$ computing RD

So there is ^{is} a monotonic lfp.

Hence RD : lfp f

Monotonic does not say

$$X \subseteq f(X) .$$

but says that

$$X \subseteq Y \Rightarrow f(X) \subseteq f(Y)$$